



Operating Systems

Lecture 2.1 - Processes and Threads

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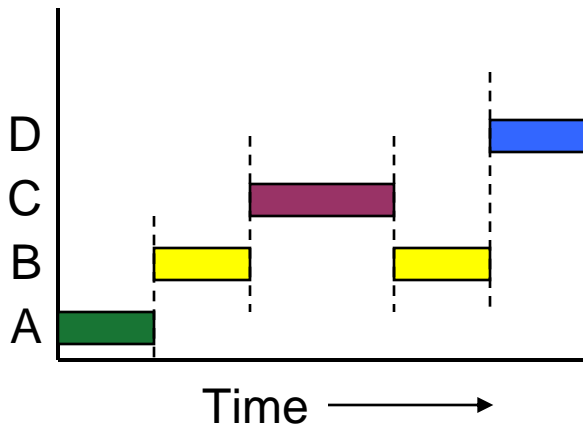
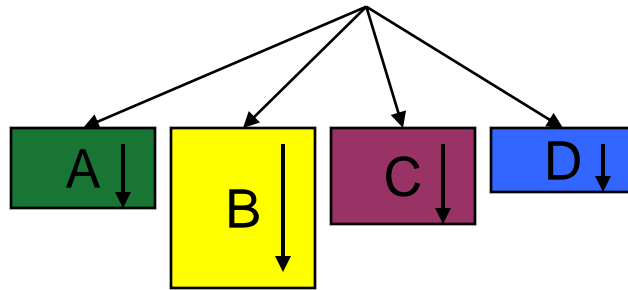
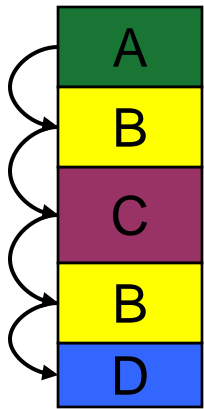
What is a process?

- Process= Program, Input/Output and State
- Program state
 - Program counter (current location in the code)
 - CPU registers
 - Stack pointer
- Only one process can be running in the CPU at any given time!
 - All executable programs are organized to multi **ordinal process**
 - Each process have a **virtual CPU (CPU switch)**

The process model

Single PC
(CPU's point of view)

Multiple PCs
(process point of view)



- **Multiprogramming** of four programs
- Conceptual model
 - 4 **independent** processes
 - Processes run **sequentially**
- Only one program active at any instant!
 - That instant can be very short...



When is a process created?

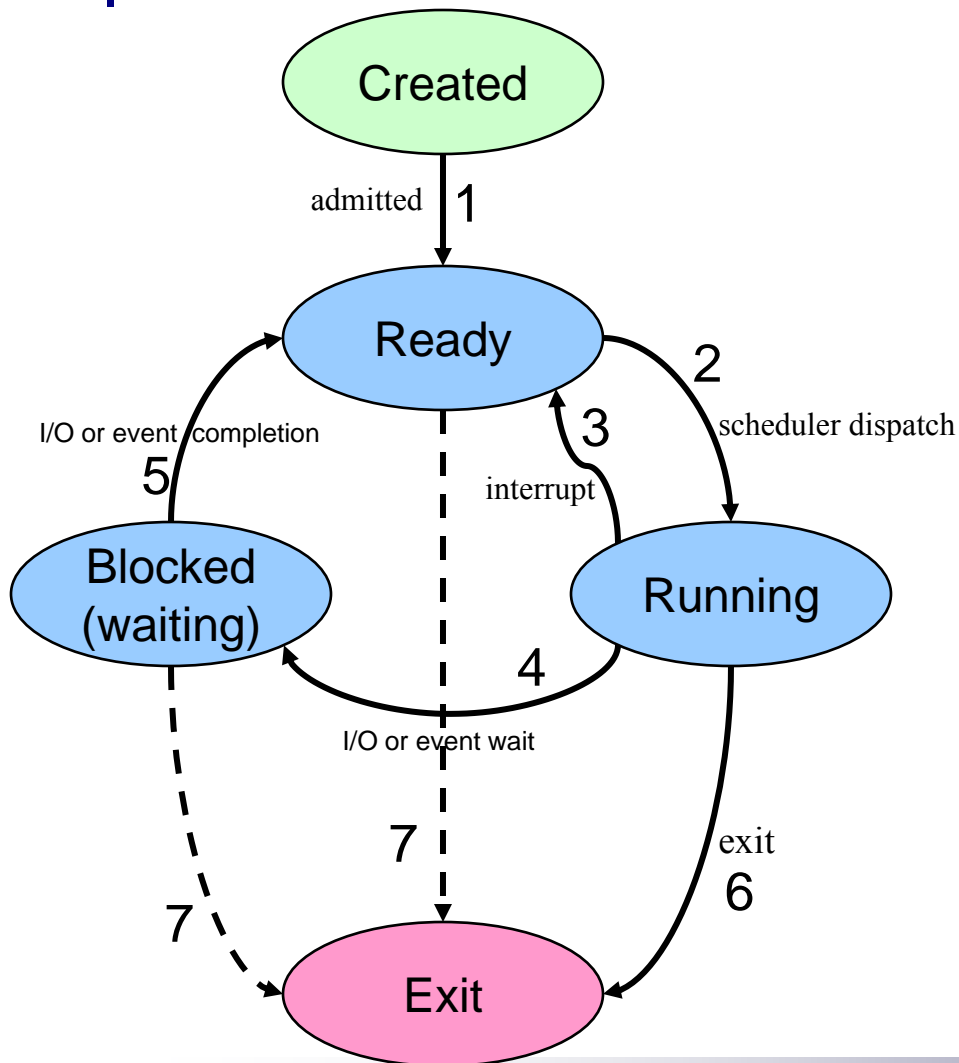
- Processes can be created in two ways
 - System initialization: one or more processes created when the **OS starts up**
 - Execution of a process **creation system call**
- System calls can come from
 - **User request** to create a new process
 - Execute a batch job



When do processes end?

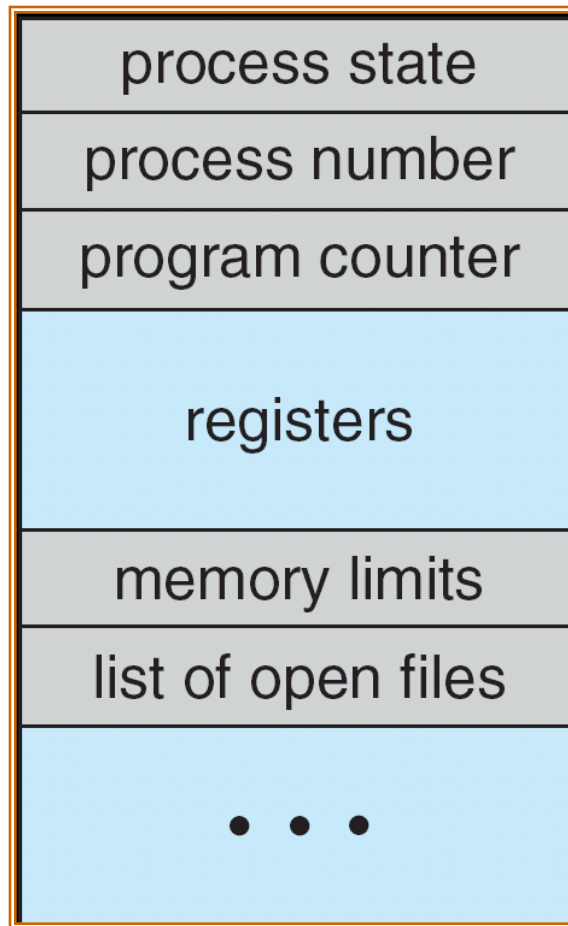
- Conditions that terminate processes can be
 - Voluntary
 - Normal exit
 - Error exit
 - Involuntary
 - Fatal error
 - Killed by another process

Process states

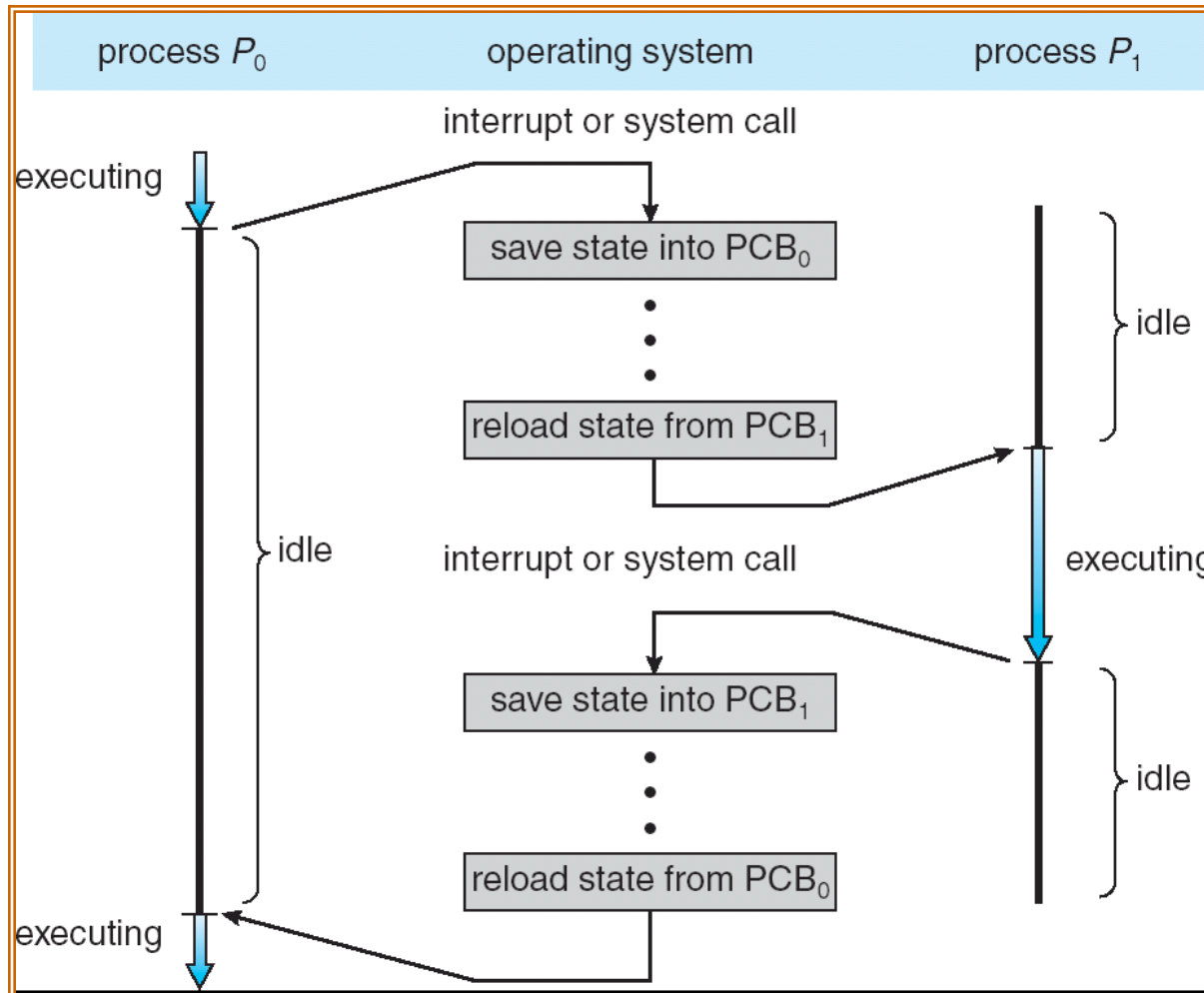


- Process in one of 5 states
 - Created
 - Ready
 - Running
 - Blocked
 - Exit
- Transitions between states
 - 1 - Process enters ready queue
 - 2 - Scheduler picks this process
 - 3 - Scheduler picks a different process
 - 4 - Process waits for event (such as I/O)
 - 5 - Event occurs
 - 6 - Process exits
 - 7 - Process ended by another process

Process Implementation: Process Control Block (PCB)



CPU Switch From Process to Process

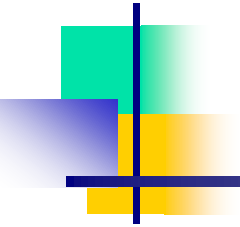




Process Scheduling Queues

- **Job queue** – set of all processes in the system
- **Ready queue** – set of all processes residing in main memory, ready and waiting to execute
- **Device queues** – set of processes waiting for an I/O device
- Processes migrate among the various queues

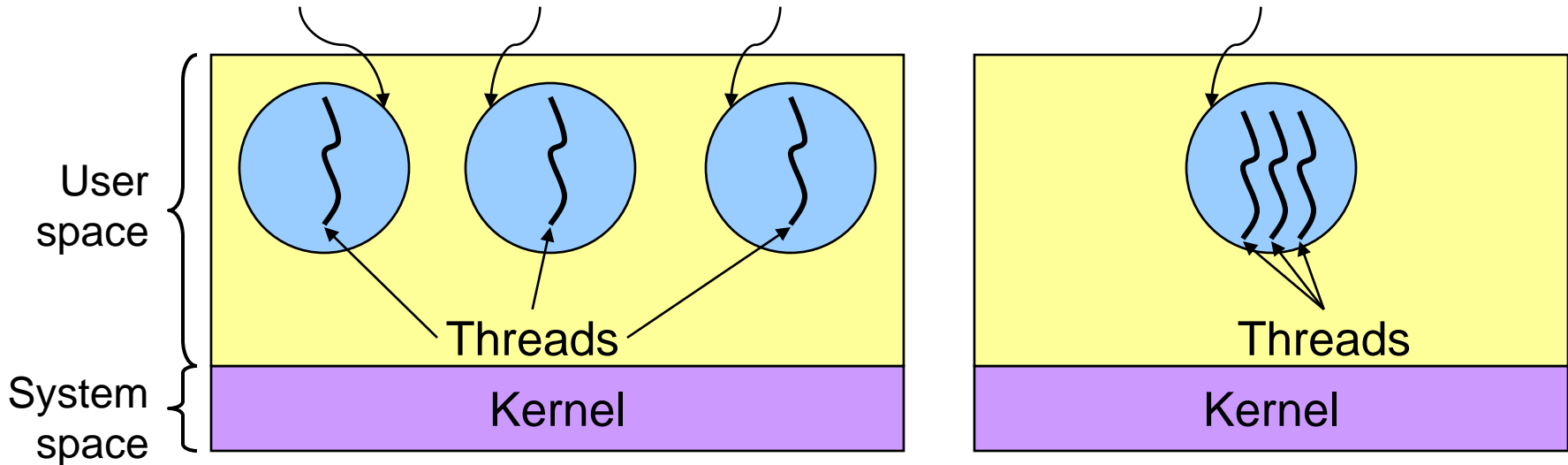
Threads



Threads: “processes” sharing memory

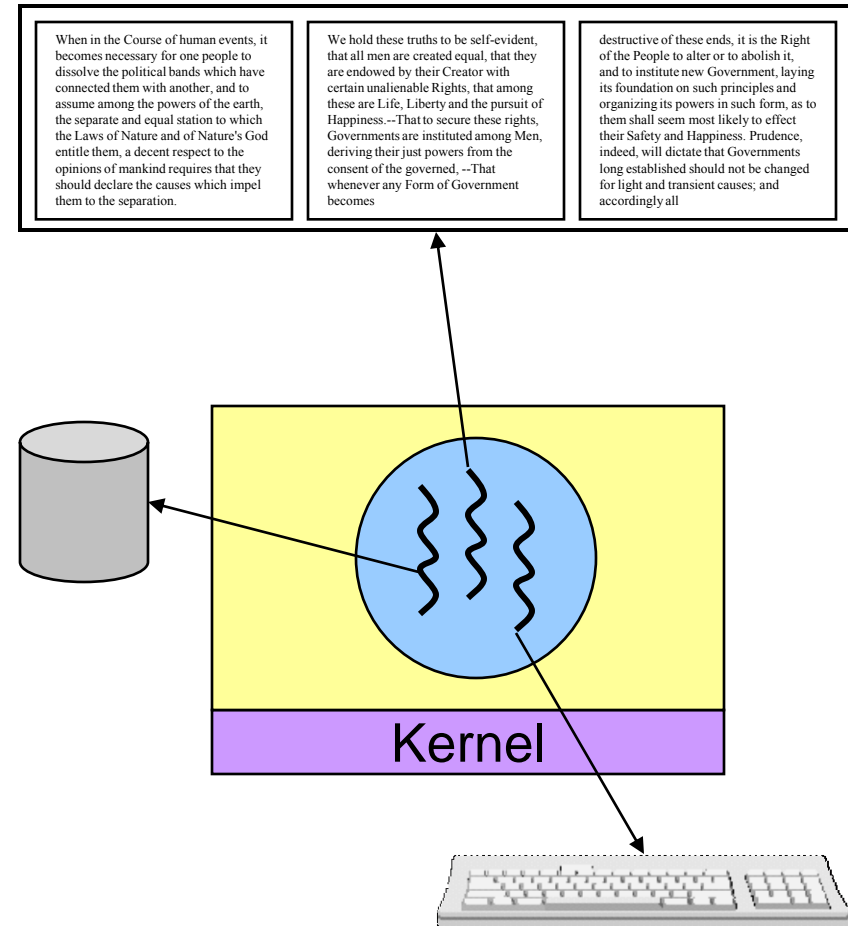
- Process == address space
- Thread == program counter / stream of instructions
- Thread == Lightweight process
- Two examples
 - Three processes, each with one thread
 - One process with three threads

Process 1 Process 2 Process 3

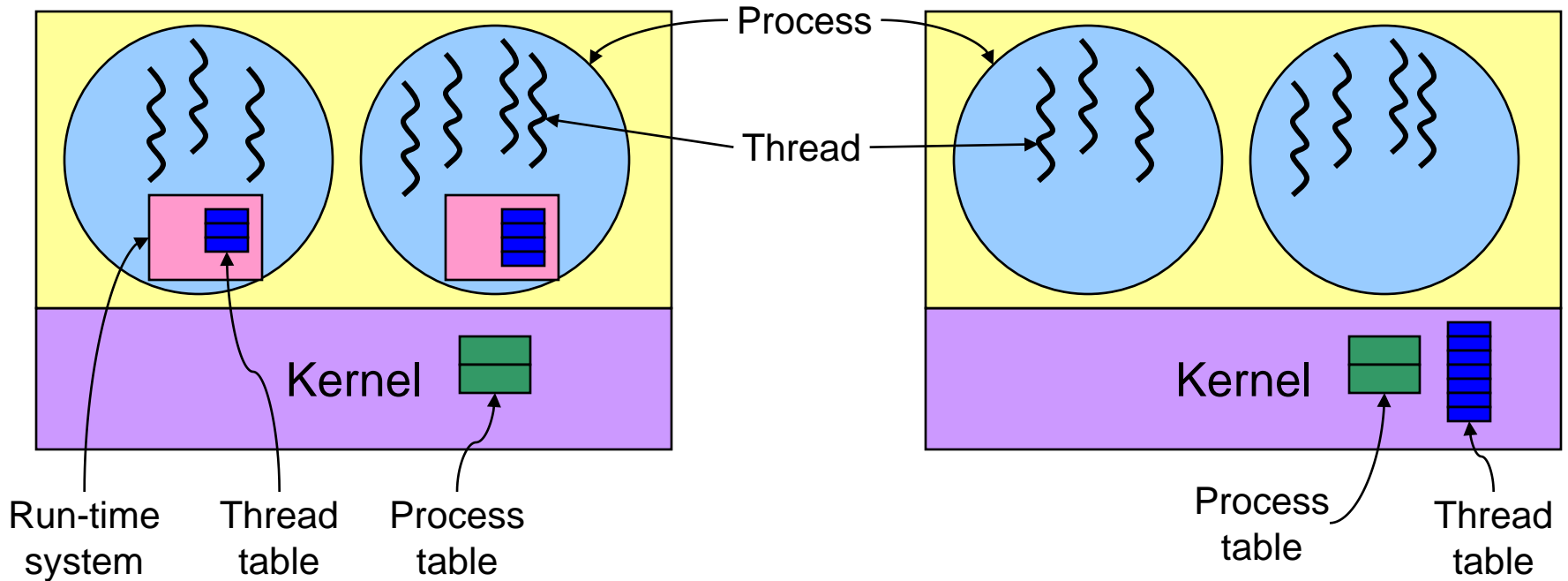


Why use threads?

- Simpler programming model
- Less waiting
- Threads are faster to create or destroy
 - No separate address space
- Overlap computation and I/O
- Example: word processor
 - Thread to **read** from **keyboard**
 - Thread to **format document**
 - Thread to **write** to **disk**



Implementing threads



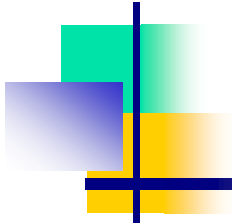
User-level threads

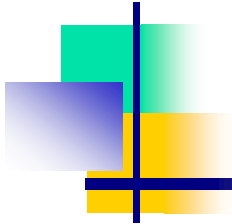
- + No need for kernel support
- May be slower than kernel threads
- Harder to do non-blocking I/O
 - Page faults

- System calls

Kernel-level threads

- + More flexible scheduling
- + Non-blocking I/O
- Not portable

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- Single Threading Vs. Multi Threading
 - Thread is a chain that include sum of the instructions
 - Address space and **some** resources is shared:
 - Open files
 - Signal handlers
 - ...
 - Each thread has own **Program counter** and **registers** and **stack**

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- Theads's Goals:
 - Increase speed of execution
 - Improve performance (Data transmission between thread of a process)
 - Parallelism
 - Modularity and Granularity