

Three Address Code Generation for Control Statements

RobinAnil (04CS3005)

- INTRODUCTION :-

- The basic idea of converting any flow of control statement to a three address code is to simulate the “branching” of the flow of control.
- This is done by skipping to different parts of the code (label) to mimic the different flow of control branches.
- Flow of control statements may be converted to three address code by use of the following functions:-
 - ✓ newlabel – returns a new symbolic label each time it is called.
 - ✓ gen () – “generates” the code (string) passed as a parameter to it.
- The following attributes are associated with the non-terminals for the code generation:-
 - ✓ code – contains the generated three address code.
 - ✓ true – contains the label to which a jump takes place if the Boolean expression associated (if any) evaluates to “true”.
 - ✓ false – contains the label to which a jump takes place if the Boolean expression (if any) associated evaluates to “false”.
 - ✓ begin – contains the label / address pointing to the beginning of the code chunk for the statement “generated” (if any) by the non-terminal.

- EXAMPLES:-

Lets try converting the following c code

FOR LOOP	in 3 TA code
<pre>a=3; b=4; for(i=0;i<n;i++){ a=b+1; a=a*a; } c=a;</pre>	<pre>a=3; b=4; i=0; L1: VAR1=i<n; if(VAR1) goto L2; goto L3; L4: i++; goto L1; L2: VAR2=b+1; a=VAR2;</pre>

	<pre> VAR3=a*a; a=VAR3; goto L4 L3: c=a; </pre>
WHILE Loop	<p>in 3 TA code</p> <pre> a=3; b=4; i=0; while(i<n){ a=b+1; a=a*a; i++; } c=a; a=3; b=4; i=0; L1: VAR1=i<n; if(VAR1) goto L2; goto L3; L2: VAR2=b+1; a=VAR2; VAR3=a*a; a=VAR3; i++; goto L1 L3: c=a; </pre>
DO WHILE Loop	<p>in 3 TA code</p> <pre> a=3; b=4; i=0; do{ a=b+1; a=a*a; i++; }while(i<n); c=a; a=3; b=4; i=0; L1: VAR2=b+1; a=VAR2; VAR3=a*a; a=VAR3; i++; VAR1=i<n; if(VAR1) goto L1; goto L2; L2: c=a; </pre>

Suppose we have the following grammar:-

$S \rightarrow \text{if } E \text{ then } S_1$
 $S \rightarrow \text{if } E \text{ then } S_1 \text{ else } S_2$
 $S \rightarrow \text{while } E \text{ do } S_1$

- SEMANTIC RULES:-

- $S \rightarrow \text{if } E \text{ then } S_1$
{
 E.true := newlabel ;
 E.false := S.next ;
 S₁.next := S.next ;
 S.code := E.code || gen(E.true ':') || S₁.code
}
 - $S \rightarrow \text{if } E \text{ then } S_1 \text{ else } S_2$
{
 E.true := newlabel ;
 E.false := newlabel ;
 S₁.next := S.next ;
 S₂.next := S.next ;
 S.code := E.code || gen(E.true ':') || S₁.code || gen('goto' S.next) || gen(E.false ':') || S₂.code
}
 - $S \rightarrow \text{while } E \text{ do } S_1$
{
 S.begin := newlabel ;
 E.true := newlabel ;
 E.false := S.next ;
 S₁.next := S.begin ;
 S.code := gen(S.begin ':') || E.code || gen(E.true ':') || S₁.code || gen('goto' S.begin)
}
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