

In the name of God

Network Flows

1. Introduction

1.2 Notation and Definitions

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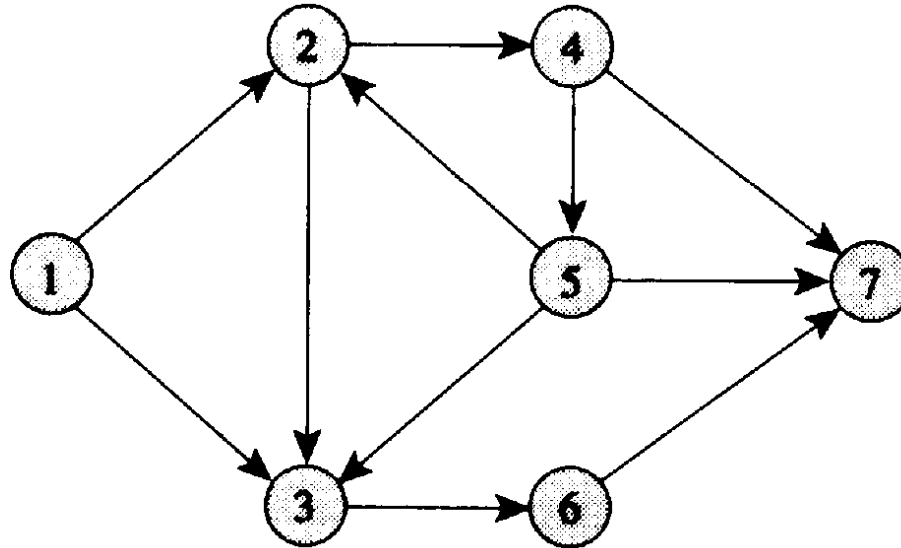
Notation and Definitions

- In this section we give several basic definitions from graph theory and present some basic notation.
- We also state some elementary properties of graphs.

Notation and Definitions

- *Directed graph*

- A directed graph $G = (N, A)$ consists of a set N of nodes and a set A of arcs whose elements are ordered pairs of distinct nodes.



- $N = \{1, 2, 3, 4, 5, 6, 7\}$
- $A = \{(1, 2), (1, 3), (2, 3), (2,4), (3, 6), (4, 5), (4, 7), (5, 2), (5, 3), (5, 7), (6, 7)\}$.

Notation and Definitions

- *Directed network*

- A *directed network* is a **directed graph** whose nodes and/or arcs have associated:

- ◆ costs
 - ◆ capacities
 - ◆ and/or supplies and demands

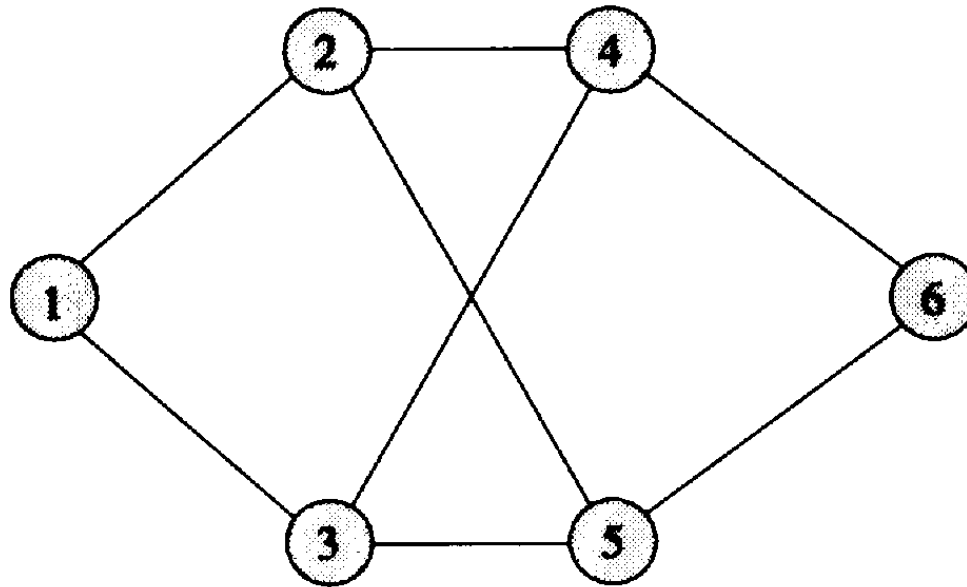
- We use the terms “graph” and “network” synonymously.

- We let n denote the number of nodes and m denote the number of arcs in G .

Notation and Definitions

- *Undirected Graphs*

- In an *undirected graph* arcs are unordered pairs of distinct nodes.



- In an undirected graph, we can refer to an arc joining the node pair i and j as either (i, j) or (j, i) .

Notation and Definitions

- *Tails and Heads*

- A directed arc (i, j) has two *endpoints* i and j .
- Node i is the *tail* of arc (i, j) and node j is its *head*.
- The arc (i, j) *emanates* from node i and *terminates* at node j .
- An arc (i, j) is *incident to* nodes i and j .
- The arc (i, j) is an *outgoing arc* of node i and an *incoming arc* of node j .
- Whenever an arc $(i, j) \in A$, we say that node j is *adjacent* to node i .

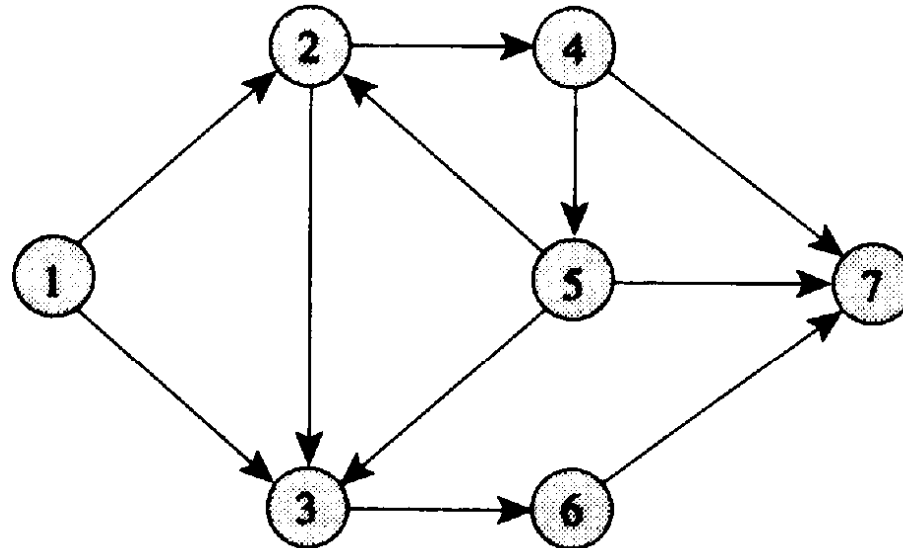
Notation and Definitions

- *Degrees*

- The *indegree* of a node is the number of **incoming arcs** of that node and its *outdegree* is the number of its **outgoing arcs**.
- The *degree* of a node is the sum of its **indegree** and **outdegree**.
- The sum of indegrees of all nodes equals the sum of outdegrees of all nodes and both are equal to the number of arcs m in the network.

Notation and Definitions

- **Example of degree:**



- Node 3 has an indegree of 3, an outdegree of 1, and a degree of 4.

Notation and Definitions

- *Adjacency List*

- The *arc adjacency list* $A(i)$ of a node i is the set of arcs emanating from that node, $A(i) = \{(i, j) \in A: j \in N\}$.
- The *node adjacency list* $A(i)$ is the set of nodes adjacent to that node; in this case, $A(i) = \{j \in N: (i, j) \in A\}$.
- We assume that arcs in the adjacency list $A(i)$ are arranged so that the head nodes of arcs are in increasing order.

Notation and Definitions

- *Multiarcs and Loops*

- *Multiarcs* are two or more arcs with the same tail and head nodes.
- A *loop* is an arc whose tail node is the same as its head node.

Notation and Definitions

- ***Subgraph***

- A graph $G' = (N', A')$ is a ***subgraph*** of $G = (N, A)$ if $N' \subseteq N$ and $A' \subseteq A$.
- We say that $G' = (N', A')$ is the subgraph of G ***induced*** by N' if A' contains each arc of A with both **endpoints** in N' .
- A graph $G' = (N', A')$ is a ***spanning subgraph*** of $G = (N, A)$ if $N' = N$ and $A' \subseteq A$.

Notation and Definitions

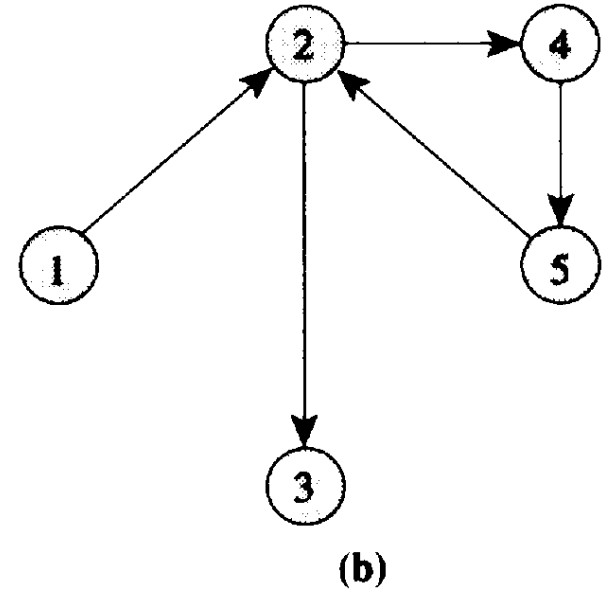
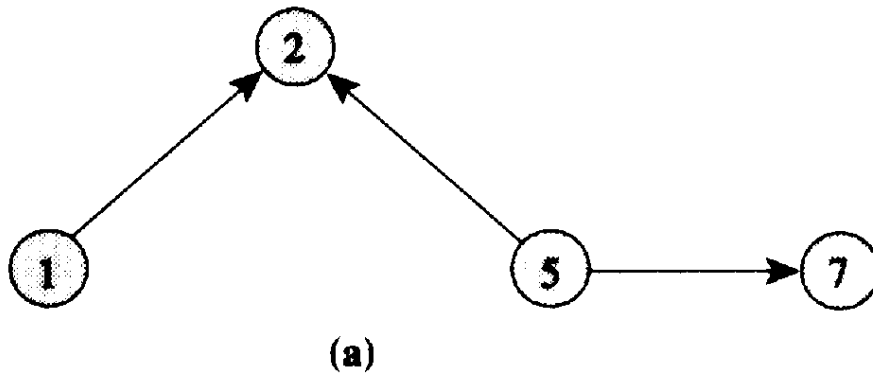
- **Walk**

- A **walk** in a directed graph $G = (N, A)$ is a subgraph of G consisting of a sequence of nodes and arcs

$$i_1 - a_1 - i_2 - a_2 - \dots - i_{r-1} - a_{r-1} - i_r$$

- satisfying the property that for all $1 \leq k \leq r - 1$,
- either $a_k = (i_k, i_{k+1}) \in A$ or $a_k = (i_{k+1}, i_k) \in A$.
- Alternatively, we shall sometimes refer to a walk as a set of (sequence of) arcs (or of nodes) without any explicit mention of the nodes (without explicit mention of arcs).

Notation and Definitions

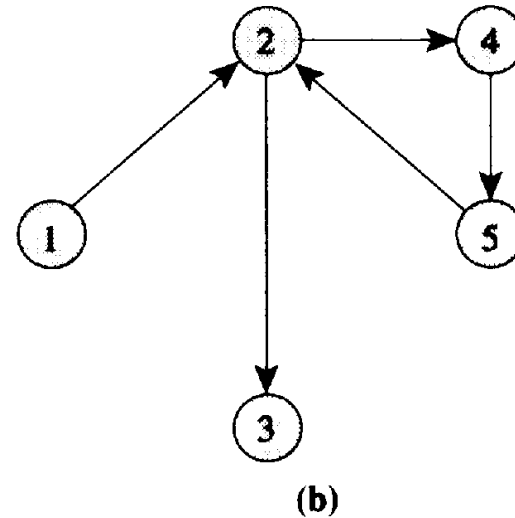
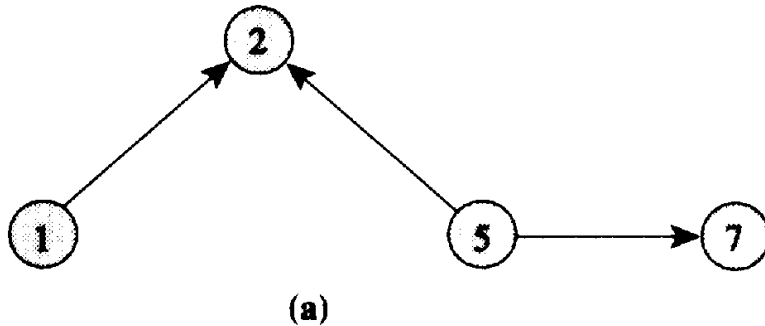


- (a) 1-2-5-7
- (b) 1-2-4-5-2-3

Notation and Definitions

- **Directed Walk**

- A **directed walk** is an oriented version of a walk in the sense that for any two consecutive nodes i_k and i_{k+1} on the walk, $(i_k, i_{k+1}) \in A$.

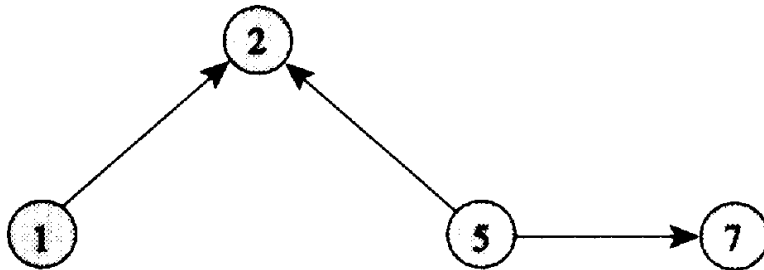


- (a) is not a directed walk
- (b) is a directed walk

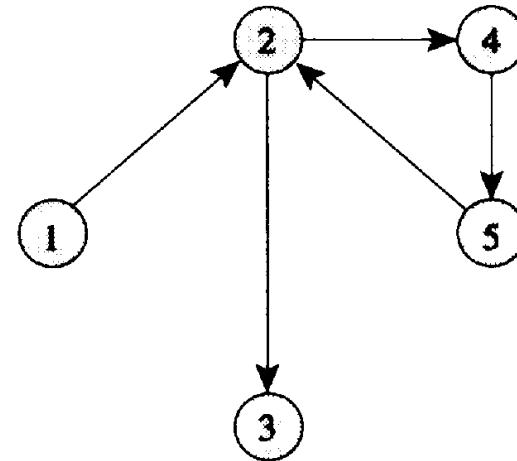
Notation and Definitions

- *Path*

- A *path* is a walk without any repetition of nodes.
- Directions are ignored.



(a)

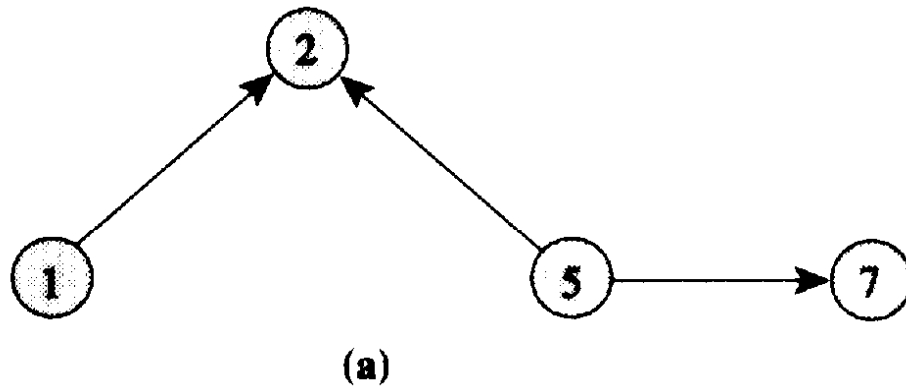


(b)

- (a) is a path, 1-2-5-7
- (b) is not a path because it repeats node 2 twice., 1-2-4-5-2-3

Notation and Definitions

- We can partition the arcs of a path into two groups: *forward arcs* and *backward arcs*.

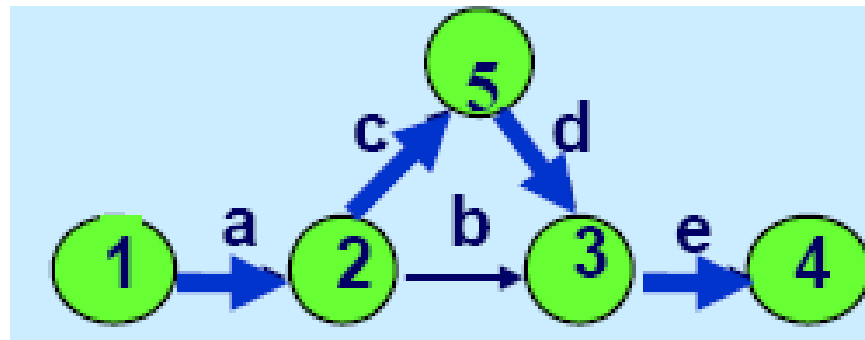


- The arcs (1, 2) and (5, 7) are forward arcs and the arc (5, 2) is a backward arc.

Notation and Definitions

- *Directed Path*

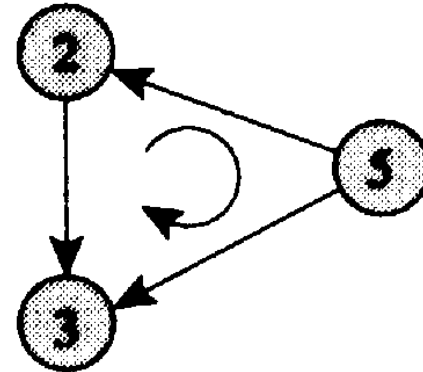
- A directed path is a directed walk without any repetition of nodes.
- In other words, a directed path has no backward arcs.
- Directions are important.



Notation and Definitions

- **Cycle**

- A **cycle** is a path $i_1 - i_2 - \dots - i_r$ together with the arc (i_r, i_1) or (i_1, i_r) .
- We refer to a cycle using the notation $i_1 - i_2 - \dots - i_r - i_1$.
- Just as we did for paths, we can define forward and backward arcs in a cycle.
- Directions are ignored.

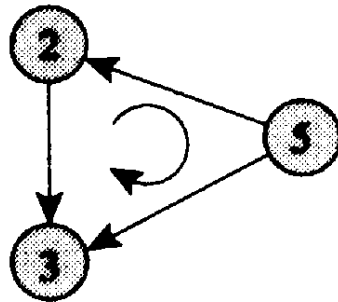


- The arcs $(5, 3)$ and $(3, 2)$ are forward arcs and the arc $(5, 2)$ is a backward arc of the cycle 2-5-3.

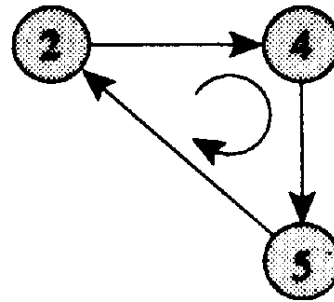
Notation and Definitions

- **Directed Cycle**

- A directed cycle is a directed path $i_1 - i_2 - \dots - i_r$ together with the arc (i_r, i_1) .
- Directions are important.



(a)



(b)

- (a) is a cycle, but not a directed cycle;
- (b) is a directed cycle.

Notation and Definitions

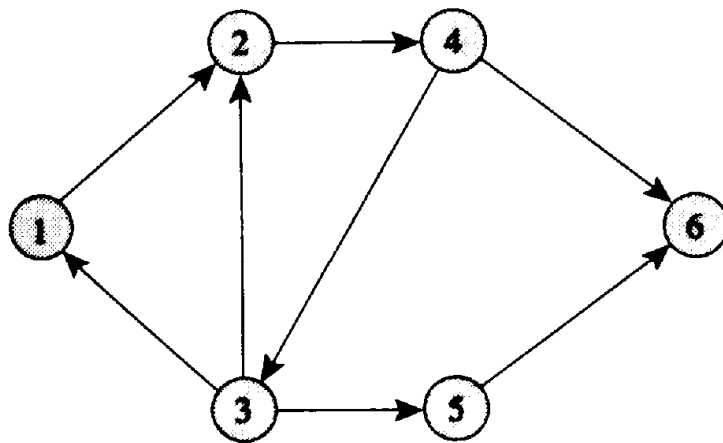
- *Acyclic Graph*

- A graph is a *acyclic* if it contains no directed cycle.

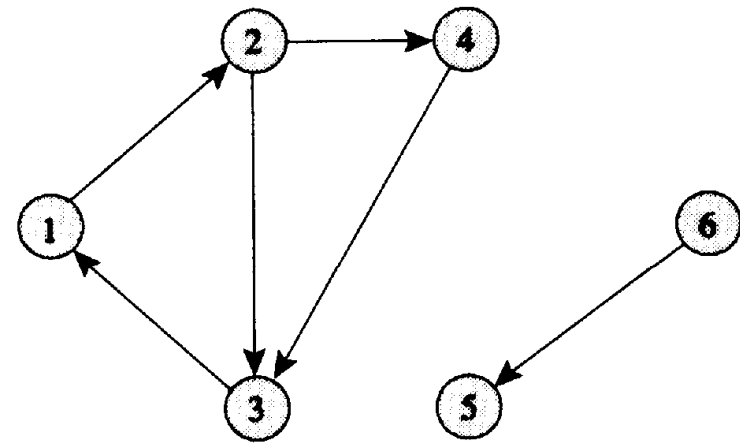
Notation and Definitions

- *Connectivity*

- We will say that two nodes i and j are *connected* if the graph contains at least **one path** from node i to node j .
- A graph is connected if every pair of its nodes is connected; otherwise, the graph is *disconnected*.



(a)



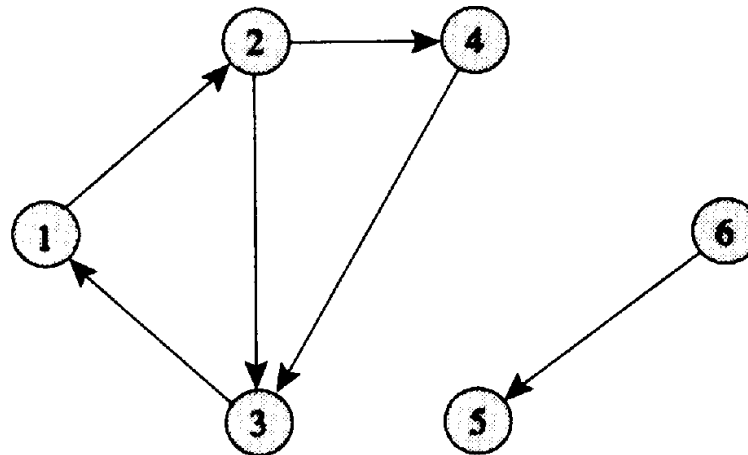
(b)

- (a) is connected graph, and (b) is disconnected graph

Notation and Definitions

- **Components**

- We refer to the maximal connected subgraphs of a disconnected network as its **components**.

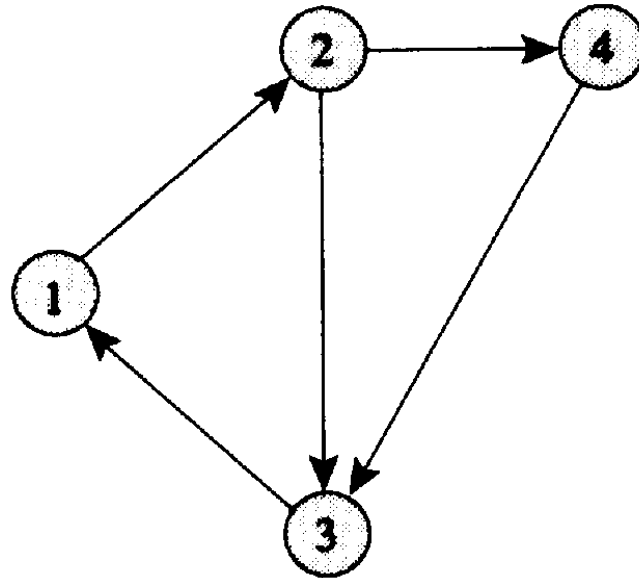


- This graph has two components consisting of the node sets $\{1, 2, 3, 4\}$ and $\{5, 6\}$.

Notation and Definitions

- ***Strong Connectivity***

- A connected graph is ***strongly connected*** if it contains at least **one directed path** from every node to every other node.



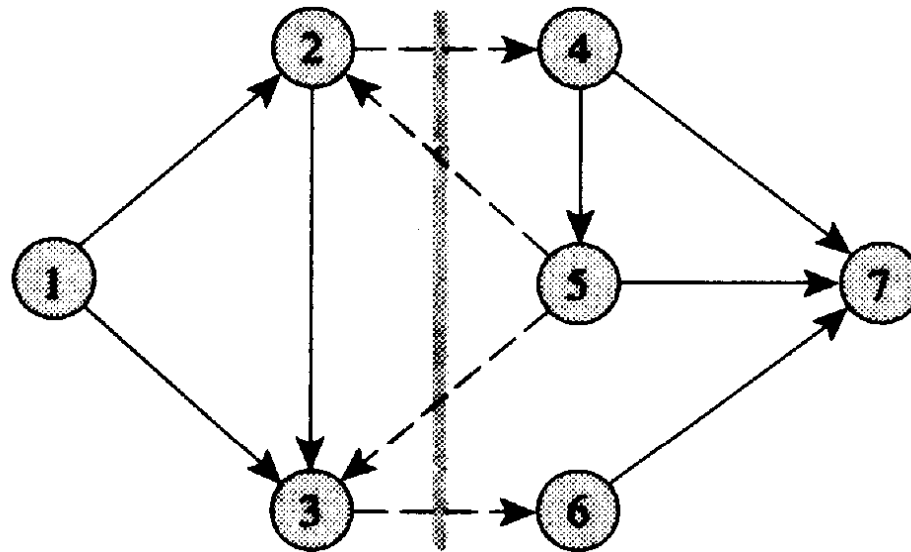
- A strongly connected graph

Notation and Definitions

- *Cut*

- A *cut* is a partition of the node set N into two parts, S and $\bar{S} = N - S$.
- Each cut defines a set of arcs consisting of those arcs that have one endpoint in S and another endpoint in \bar{S} .
- We refer to this set of arcs as a cut and represent it by the notation $[S, \bar{S}]$.

Notation and Definitions

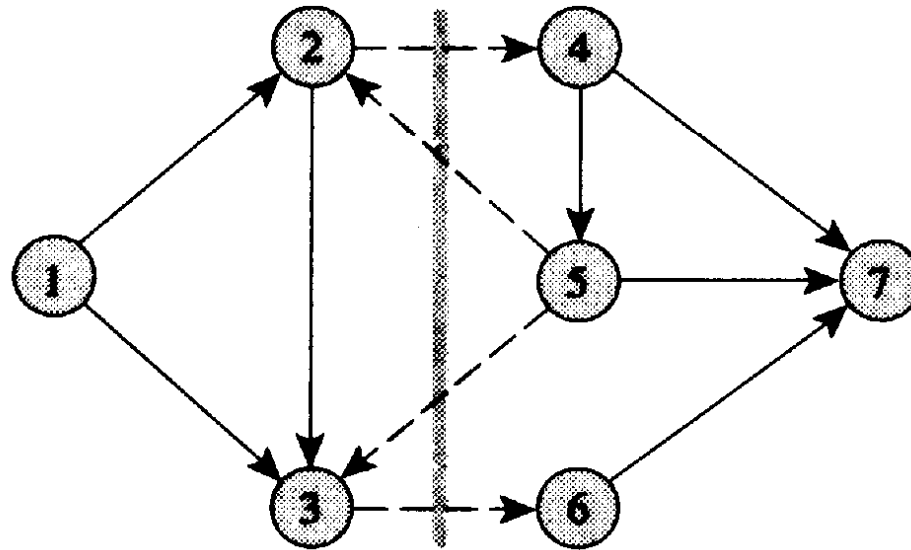


- A cut with $S = \{1, 2, 3\}$ and $\bar{S} = \{4, 5, 6, 7\}$.
- The set of arcs in this cut are $\{(2, 4), (5, 2), (5, 3), (3, 6)\}$.

Notation and Definitions

- *s-t Cut*

- An *s-t* cut is defined with respect to two distinguished nodes s and t , and is a cut $[S, \bar{S}]$ satisfying the property that $s \in S$ and $t \in \bar{S}$.



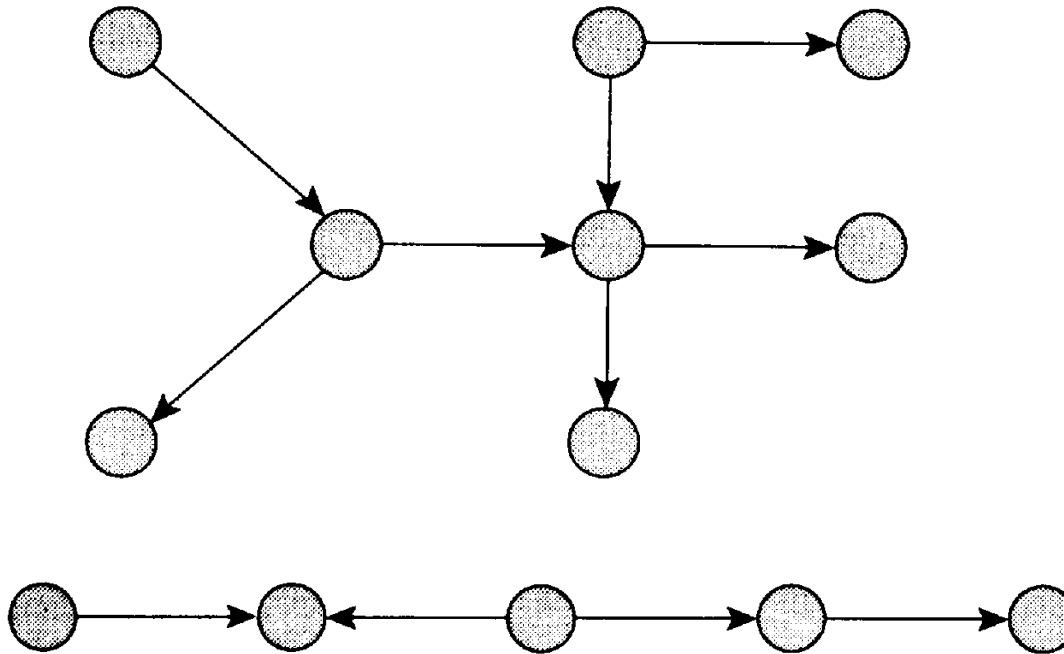
- If $s = 1$ and $t = 6$, this cut is an *s-t* cut;
- but if $s = 1$ and $t = 3$, this cut is not an *s-t* cut.

Notation and Definitions

- *Tree*
 - A *tree* is a connected graph that contains no cycle.
- Property
 - (a) A tree on n nodes contains exactly $n - 1$ arcs.
 - (b) A tree has at least two leaf nodes (i.e., nodes with degree 1).
 - (c) Every two nodes of a tree are connected by a unique path.

Notation and Definitions

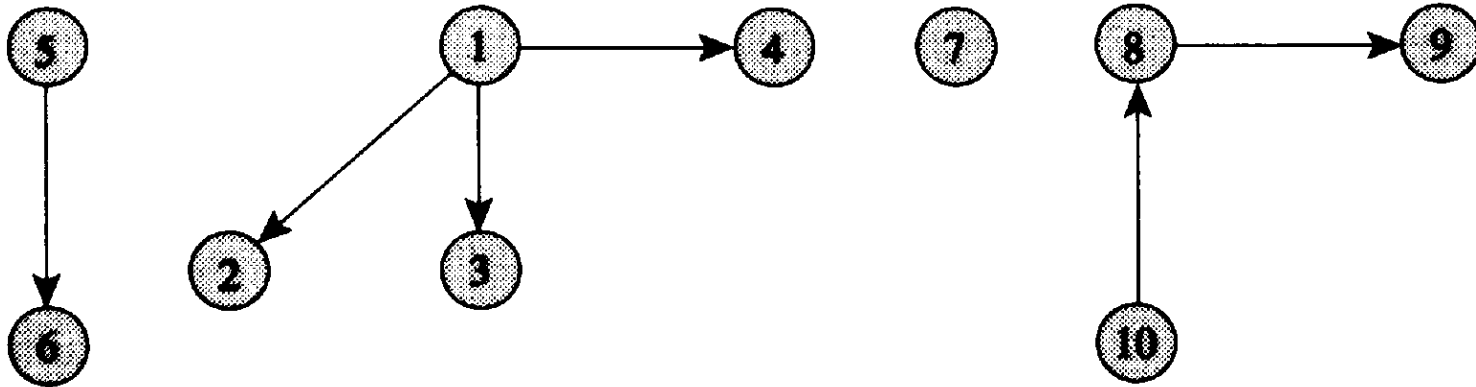
- Two examples of trees



Notation and Definitions

- *Forest*

- A graph that contains no cycle is a *forest*.
- Alternatively, a forest is a collection of trees.



- an example of a forest.

Notation and Definitions

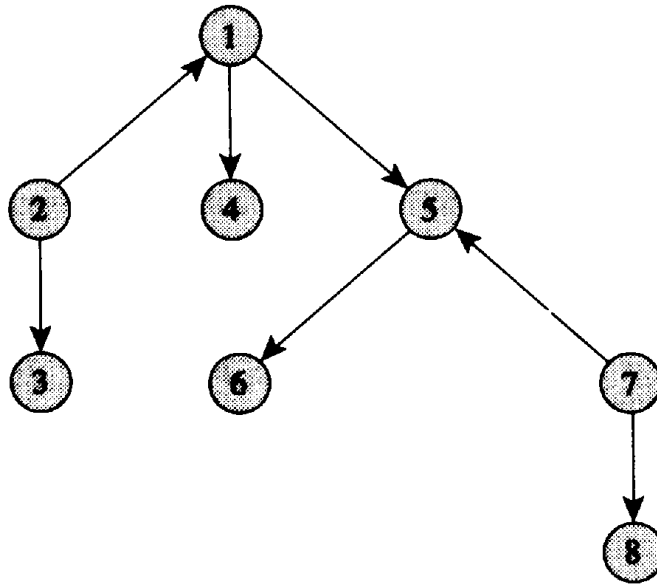
- *Subtree*

- A connected subgraph of a tree is a *subtree*.

Notation and Definitions

- ***Rooted Tree***

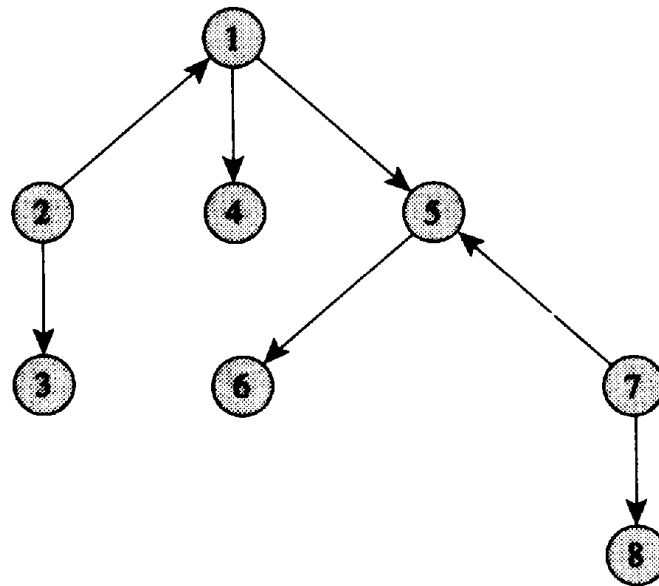
- A rooted tree is a tree with a specially designated node, called its ***root***; we regard a rooted tree as though it were hanging from its root.



- an instance of a rooted tree; node 1 is the root node.

Notation and Definitions

- We often view the arcs in a rooted tree as defining predecessor-successor (or parent-child) relationships.



- Node 5 is the predecessor of nodes 6 and 7,
- Node 1 is the predecessor of nodes 2, 4, and 5.

Notation and Definitions

- Each node i (except the root node) has a unique predecessor, which is the next node on the unique path in the tree from that node to the root;
- We store the predecessor of node i using a predecessor index $pred(i)$.
- If $j = pred(i)$, we say that node j is the predecessor of node i and node i is a successor of node j .
- These predecessor indices uniquely define a rooted tree and also allow us to trace out the unique path from any node back to the root.

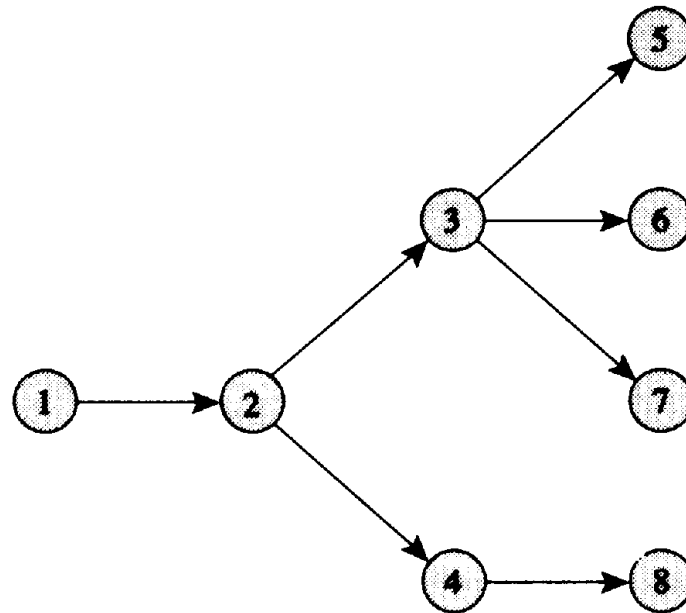
Notation and Definitions

- There are two special type of rooted trees, called a *directed in-tree* and a *directed out-tree*.

Notation and Definitions

- *Directed-Out- Tree*

- A tree is a *directed out-tree* rooted at node s if the unique path in the tree from node s to every other node is a directed path.

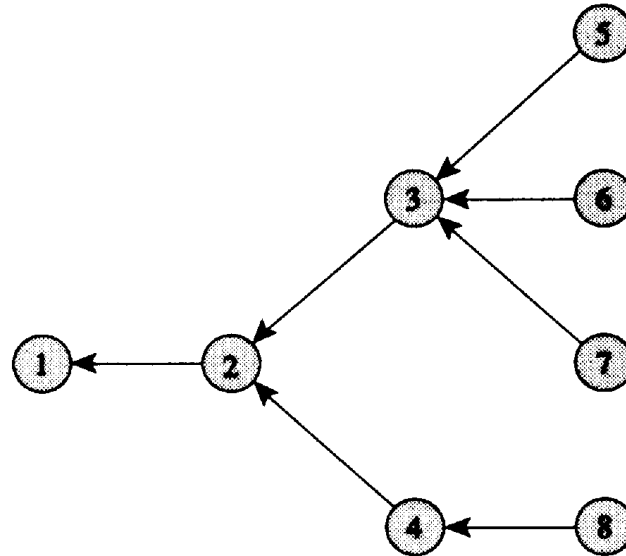


- an instance of a directed out-tree rooted at node 1.
- Observe that every node in the directed outtree (except node 1) has indegree 1.

Notation and Definitions

- *Directed-In-Tree*

- A tree is a *directed in-tree* rooted at node s if the unique path in the tree from any node to node s is a directed path.



- an instance of a directed in-tree rooted at node 1.
- Observe that every node in the directed in-tree (except node 1) has outdegree 1.

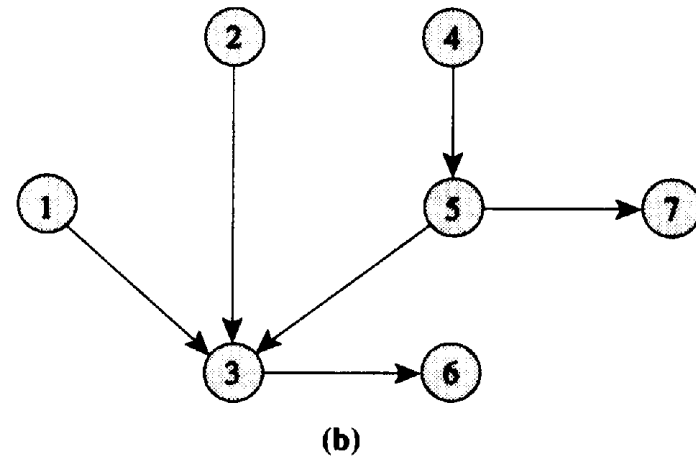
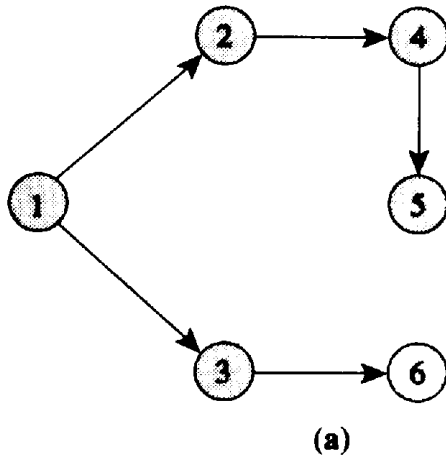
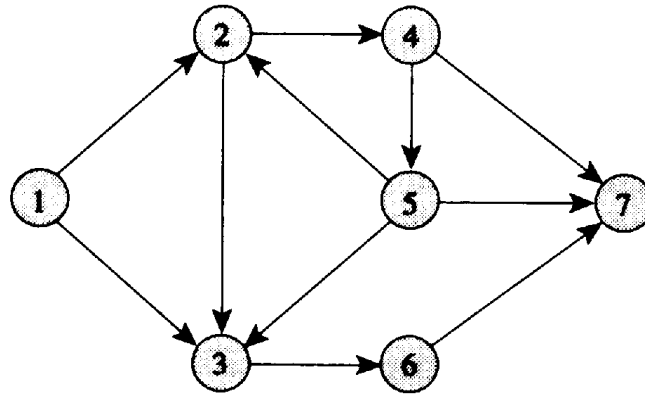
Notation and Definitions

- *Spanning Tree*

- A tree T is a spanning tree of G if T is a spanning subgraph of G .
- Every spanning tree of a connected n -node graph G has $(n - 1)$ arcs.
- We refer to the arcs belonging to a spanning tree T as tree arcs and arcs not belonging to T as nontree arcs.

Notation and Definitions

- two spanning trees of the graph



Notation and Definitions

- ***Undirected networks***

- An undirected arc (i, j) has two endpoints, i and j , but its tail and head nodes are undefined.
- If the network contains the arc (i, j) , node i is adjacent to node j , and node j is adjacent to node i .



The End