

**In the name of God**

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# **Network Flows**

## **6. Lagrangian Relaxation**

### **6.2 Solving the Lagrangian Multiplier Problem**

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# Outline

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- Introduction
- Subgradient Optimisation Technique
- Subgradient Optimization and Inequality Constraints

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# **Introduction**

# Solving the Lagrangian Multiplier Problem

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- Consider the *constrained shortest path problem*
- Suppose that now we have a time limitation of  $T = 14$  instead of  $T = 10$ .
- When we relax the time constraint, the *Lagrangian multiplier function*  $L(\mu)$  becomes:

$$L(\mu) = \min\{c_P + \mu(t_P - T) : P \in \mathcal{P}\}.$$

- where,  $\mathcal{P}$  is the collection of all directed paths from the source node 1 to the sink node  $n$ .

# Solving the Lagrangian Multiplier Problem

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- For convenience, we refer to the quantity  $c_p + \mu(t_p - T)$  as the *composite cost* of the path  $P$ .
- For a specific value of the Lagrangian multiplier  $\mu$ , we can solve  $L(\mu)$  by enumerating all the directed paths in  $\mathcal{P}$  and choosing the path with the smallest composite cost.
- We can solve the Lagrangian multiplier problem by determining  $L(\mu)$  for all nonnegative values of the Lagrangian multiplier  $\mu$  and choosing the value that achieves  $\max_{\mu \geq 0} L(\mu)$ .

# Solving the Lagrangian Multiplier Problem

- Path cost and time data with  $T = 14$

<b>Path <math>P</math></b>	<b>Path cost <math>c_P</math></b>	<b>Path time <math>t_P</math></b>	<b>Composite cost <math>c_P + \mu (t_P - T)</math></b>
1-2-4-6	3	18	$3 + 4\mu$
1-2-5-6	5	15	$5 + \mu$
1-2-4-5-6	14	14	14
1-3-2-4-6	13	13	$13 - \mu$
1-3-2-5-6	15	10	$15 - 4\mu$
1-3-2-4-5-6	24	9	$24 - 5\mu$
1-3-4-6	16	17	$16 + 3\mu$
1-3-4-5-6	27	13	$27 - \mu$
1-3-5-6	24	8	$24 - 6\mu$

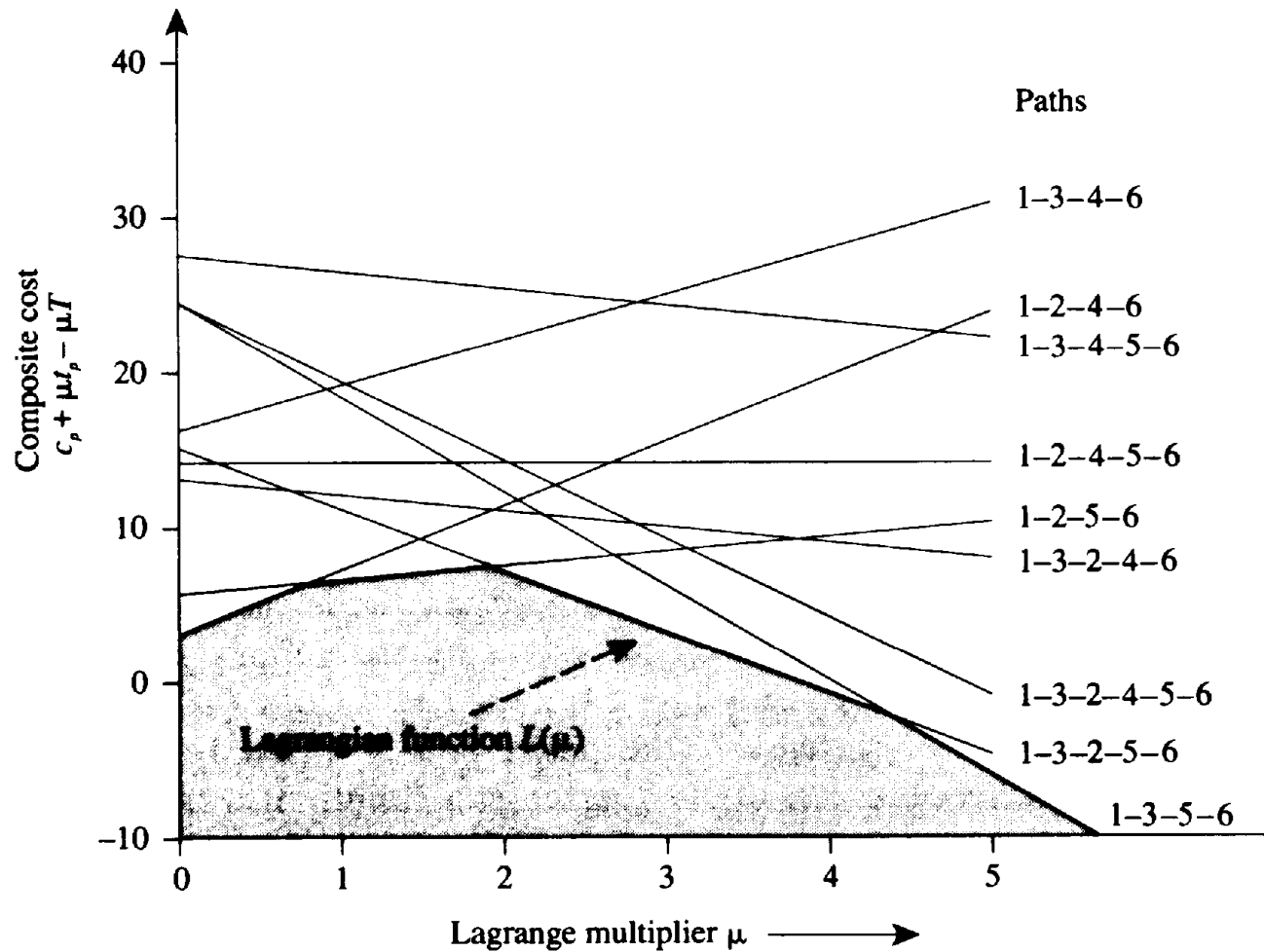
# Solving the Lagrangian Multiplier Problem

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- The composite cost  $c_P + \mu(t_P - T)$  for any path  $P$  is a linear function of  $\mu$  with an intercept of  $c_P$  and a slope of  $(t_P - T)$

# Solving the Lagrangian Multiplier Problem

- Lagrangian function,  $T = 14$





# Solving the Lagrangian Multiplier Problem

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- To find the *optimal multiplier value*  $\mu^*$  of the Lagrangian multiplier problem, we need to find *the highest point* of the *Lagrangian multiplier function*  $L(\mu)$ .
- Suppose that we consider the polyhedron defined by those points that lie on or below the function  $L(\mu)$ .
- These are the shaded points in the Figure.
- Then geometrically, we are finding the highest point in a polyhedron defined by the function  $L(\mu)$ , which is a linear program.

# Solving the Lagrangian Multiplier Problem

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- Consider the generic optimization model (P), defined

$$\min\{cx : Ax = b, x \in X\}$$

- and suppose that the set  $X$  is finite.

$$X = \{x^1, x^2, \dots, x^K\}$$

- By relaxing the constraints  $Ax = b$ , we obtain the Lagrangian multiplier function:

$$L(\mu) = \min\{cx + \mu(Ax - b) : x \in X\}.$$

- By definition:

$$L(\mu) \leq cx^k + \mu(Ax^k - b) \quad \text{for all } k = 1, 2, \dots, K.$$

# Solving the Lagrangian Multiplier Problem

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- In the space of composite costs and Lagrange multipliers  $\mu$ , each function  $cx^k + \mu(Ax^k - b)$  is a multidimensional "line" called a hyperplane (if  $\mu$  is two-dimensional, it is a plane).
- The Lagrangian multiplier function  $L(\mu)$  is the lower envelope of the hyperplanes  $cx^k + \mu(Ax^k - b)$  for  $k = 1, 2, \dots, K$ .
- In the Lagrangian multiplier problem, we wish to determine the highest point on this envelope

# Solving the Lagrangian Multiplier Problem

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- We can find this the highest point by solving the optimization problem:

Maximize  $w$

subject to

$$w \leq cx^k + \mu(Ax^k - b) \quad \text{for all } k = 1, 2, \dots, K,$$

$\mu$  unrestricted,

# Solving the Lagrangian Multiplier Problem

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- Theorem.

- The Lagrangian multiplier problem  $L^* = \max_{\mu} L(\mu)$  with

$$L(\mu) = \min\{cx^k + \mu(Ax - b) : x \in X\}$$

- is equivalent to the linear programming problem

$$L^* = \max\{w : w \leq cx^k + \mu(Ax^k - b) \\ \text{for } k = 1, 2, \dots, K\}.$$

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# **Subgradient Optimisation Technique**

# Subgradient Optimisation Technique

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- In solving optimization problems with the nonlinear objective function  $f(\mathbf{x})$  of an  $n$ -dimensional vector  $\mathbf{x}$ , researchers and practitioners often use *gradient methods*.

# Subgradient Optimisation Technique

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- Suppose that in solving the *Lagrangian multiplier problem*, we are at a point where the Lagrangian function has a unique solution  $\bar{x}$ .

$$L(\mu) = \min\{cx + \mu(Ax - b) : x \in X\}$$

- Since

$$L(\mu) = c\bar{x} + \mu(A\bar{x} - b)$$

- The solution  $\bar{x}$  remains optimal for small changes in the value of  $\mu$ , the *gradient* at this point is

$$A\bar{x} - b$$



# Subgradient Optimisation Technique

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- A gradient method would change the value of  $\mu$  as follows:  $\mu \leftarrow \mu + \theta(\mathcal{A}\bar{x} - b)$ .
  - $\theta$  : is a step size (a scalar) that specifies how far we move in the gradient direction.
  - If  $(\mathcal{A}x - b)_i = 0$ , the solution  $x$  uses up exactly the required units of the  $i$ th resource, and we hold the Lagrange multiplier (the toll)  $\mu_i$  of that resource at its current value;
  - If  $(\mathcal{A}x - b)_i < 0$ , the solution  $x$  uses up less than the available units of the  $i$ th resource and we decrease the Lagrange multiplier  $\mu_i$  on that resource;
  - If  $(\mathcal{A}x - b)_i > 0$ , the solution  $x$  uses up more than the available units of the  $i$ th resource and we increase the Lagrange multiplier  $\mu_i$  on that resource.

# Subgradient Optimisation Technique

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- To solve the Lagrangian multiplier problem, let  $\mu^0$  be any initial choice of the Lagrange multiplier;
- we determine the subsequent values  $\mu^k$  for  $k = 1, 2, \dots$ , of the Lagrange multipliers as follows:

$$\mu^{k+1} = \mu^k + \theta_k(Ax^k - b).$$

- where  $x^k$  : is any solution to the Lagrangian subproblem when  $\mu = \mu^k$  and  $\theta_k$  is the step length at the  $k$ th iteration.

# Subgradient Optimisation Technique

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- To ensure that this method solves the Lagrangian multiplier problem, we need to exercise some care in the choice of the step sizes  $\theta$ .
  - If we choose them too small, the algorithm would become stuck at the current point and not converge;
  - If we choose the step sizes too large, the iterates  $\mu^k$  might overshoot the optimal solution and perhaps even oscillate between two nonoptimal solutions.

# Subgradient Optimisation Technique

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- *Newton's method*

- It is an important variant of the subgradient optimization procedure
- Suppose that

$$L(\mu^k) = cx^k + \mu^k(Ax^k - b)$$

- that is,  $x^k$  solves the Lagrangian subproblem when  $\mu = \mu^k$
- We assume that  $x^k$  continues to solve the Lagrangian subproblem as we vary  $\mu$ ; or, stated in another way, we make a linear approximation

$$r(\mu) = cx^k + \mu(Ax^k - b) \text{ to } L(\mu)$$

# Subgradient Optimisation Technique

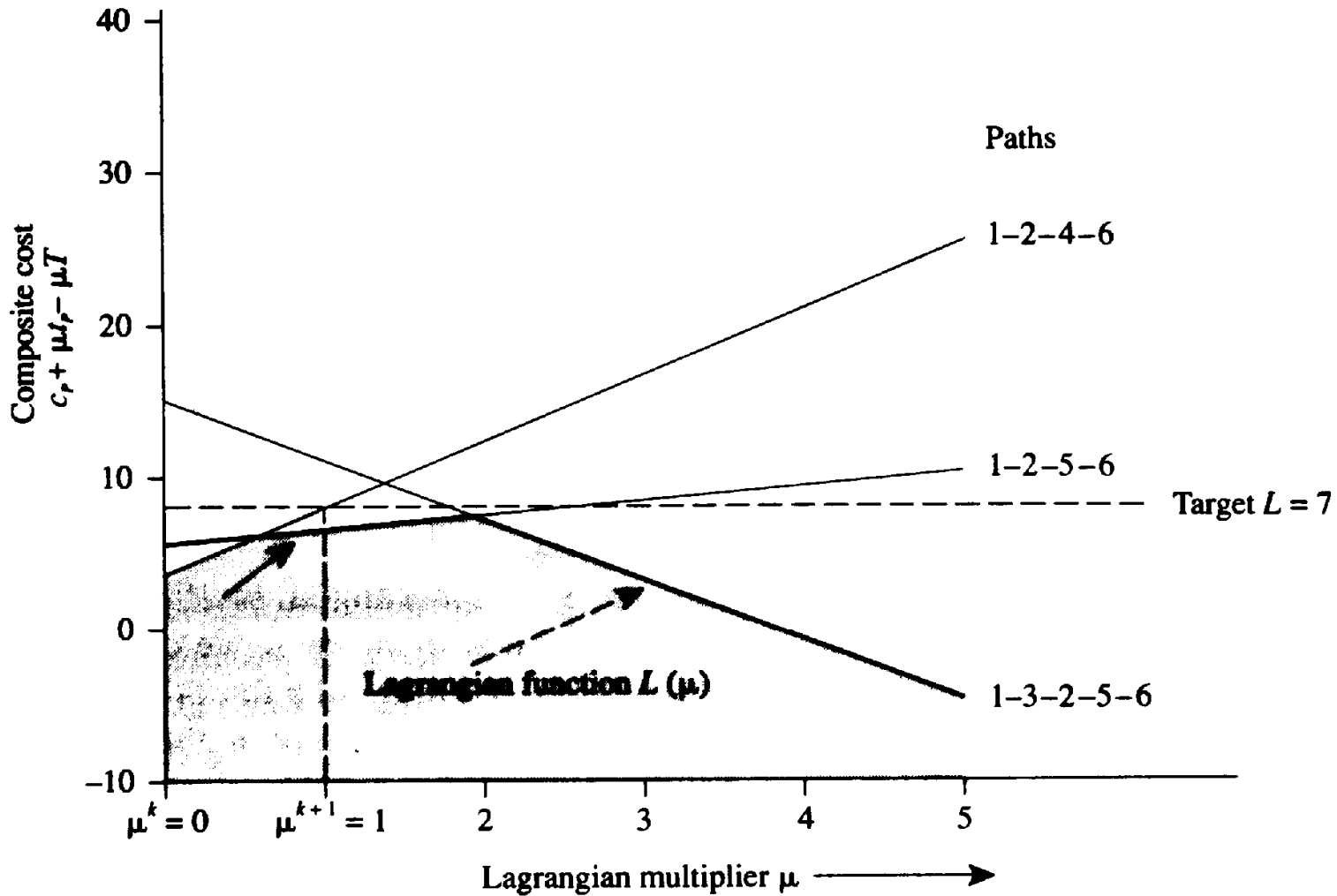
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- *Newton's method (cont.)*

- Suppose that we know the optimal value  $L^*$  of the Lagrangian multiplier problem (which we do not).
- Then we might move in the subgradient direction until the value of the linear approximation exactly equals  $L^*$ .

# Subgradient Optimisation Technique

- The constrained shortest path example,  $T=14$



# Subgradient Optimisation Technique

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- **The constrained shortest path example**

- The figure shows an example of Newton's method when applied to the constrained shortest path example, starting with  $\mu^k = 0$ .
- At this point, the path  $P = 1-2-4-6$  solves the Lagrangian subproblem and  $\mathcal{A}x - \mathbf{b}$  equals  $t_p - \mathbf{T} = 18 - 14 = 4$ .
- Since  $L^* = 7$  and the path  $P$  has a cost  $c_p = 3$ , in accordance with this linear approximation, or Newton's method, we would approximate

$$L(\mu) \text{ by } r(\mu) = 3 + 4\mu, \text{ set } 3 + 4\mu = 7$$

- and define the new value of  $\mu$  as  
$$\mu^{k+1} = (7 - 3)/4 = 1$$

# Subgradient Optimisation Technique

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- In general, we set the step length  $\theta_k$  so that

$$r(\mu^{k+1}) = cx^k + \mu^{k+1}(Ax^k - b) = L^*,$$

- since,

$$\mu^{k+1} = \mu^k + \theta_k(Ax^k - b),$$

- then,

$$r(\mu^{k+1}) = cx^k + [\mu^k + \theta_k(Ax^k - b)](Ax^k - b) = L^*.$$



# Subgradient Optimisation Technique

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- recalling that

$$L(\mu^k) = cx^k + \mu(\mathcal{A}x^k - b)$$

- and letting the Euclidean norm of the vector  $y$ :

$$\|y\| = \left(\sum_j y_j^2\right)^{1/2}$$

- we can solve for the *step length* and find that

$$\theta_k = \frac{L^* - L(\mu^k)}{\|\mathcal{A}x^k - b\|^2}.$$

# Subgradient Optimisation Technique

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- Since we do not know the optimal objective function value  $L^*$  of the Lagrangian multiplier problem,
  - practitioners of Lagrangian relaxation often use the following popular heuristic for selecting the *step length*:

$$\theta_k = \frac{\lambda_k [\text{UB} - L(\mu^k)]}{\|Ax^k - b\|^2}.$$

- ◆ **UB** : is an upper bound on the optimal objective function value  $z^*$  of the problem (P), and so an upper bound on  $L^*$  as well
- ◆  $\lambda_k$  : is a scalar chosen (strictly) between 0 and 2.

# Subgradient Optimisation Technique

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- *The heuristic procedure:*

- Initially, the upper bound is the objective function value of any known feasible solution to the problem (P).
- As the algorithm proceeds, if it generates a better (i.e., lower cost) feasible solution, it uses the objective function value of this solution in place of the upper bound **UB**.
- Usually, practitioners choose the scalars  $\lambda_k$  by starting with  $\lambda_k = 2$  and then reducing  $\lambda_k$  whenever the best Lagrangian objective function value found so far has failed to increase in a specified number of iterations.
- Since this version of the algorithm has no convenient stopping criteria, practitioners usually terminate it after it has performed a specified number of iterations.

# Subgradient Optimisation Technique

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- we might note that the subgradient optimization procedure is not the only way to solve the Lagrangian multiplier problem:
  - practitioners have used a number of other heuristics, including methods known as *multiplier ascent methods*.

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# **Subgradient Optimization and Inequality Constraints**

# Subgradient Optimization and Inequality Constraints

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- If we apply Lagrangian relaxation to a problem with constraints  $\mathcal{A}x - b$  stated in inequality form instead of the equality constraints, the Lagrange multipliers  $\mu$  are constrained to be nonnegative.

- The update formula

$$\mu^{k+1} = \mu^k + \theta_k(\mathcal{A}x^k - b)$$

- might cause one or more of the components  $\mu_i$  of  $\mu$  to become negative.

# Subgradient Optimization and Inequality Constraints

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- To avoid this possibility, we modify the update formula as follows:

$$\mu^{k+1} = [\mu^k + \theta_k(\mathcal{A}x^k - b)]^+$$

- where, the notation  $[y]^+$  denotes the "positive part" of the vector  $y$ ; that is, the  $i$ th component of  $[y]^+$  equals the maximum of  $0$  and  $y_i$ .
- Stated in another way, if the update formula

$$\mu^{k+1} = \mu^k + \theta_k(\mathcal{A}x^k - b)$$

- would cause the  $i$ th component of  $\mu_i$  to be negative, then we simply set the value of this component to be zero.

# Subgradient Optimization and Inequality Constraints

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- We then implement all the other steps of the subgradient procedure exactly the same as for problems with equality constraints.
  - i.e., the choice of the step size  $\rho$  at each step and
  - the solution of the Lagrangian subproblems
- For problems with both equality and inequality constraints, we use a mixture of the equality and inequality versions of the algorithm



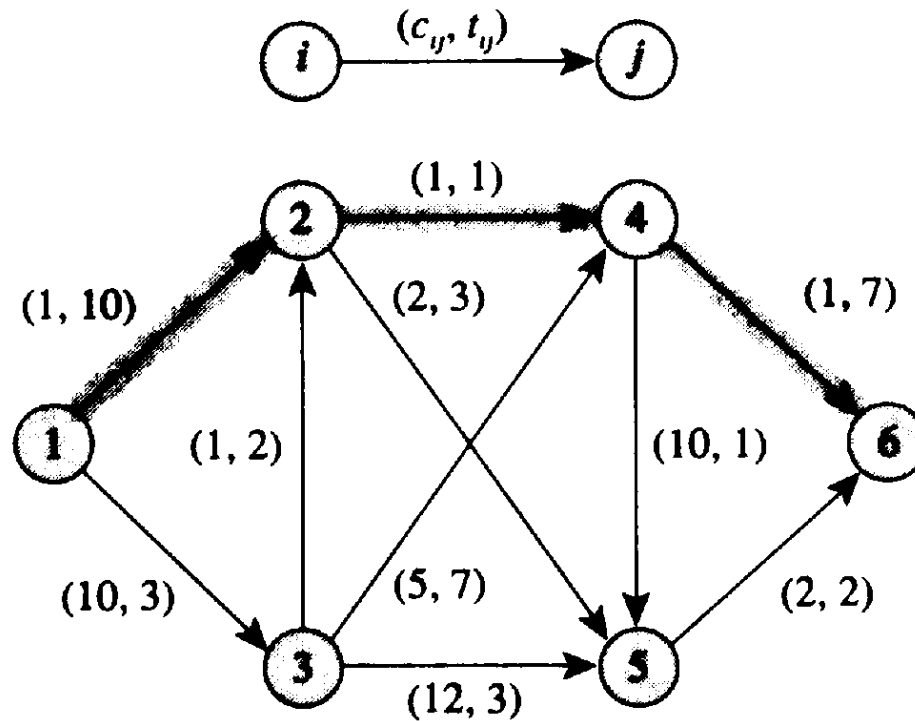
# Subgradient Optimization and Inequality Constraints

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- **The constrained shortest path example:**

- We start to solve our constrained shortest path problem at  $\mu^0 = \mathbf{0}$  with  $\lambda^0 = 0.8$  and with  $\mathbf{UB} = 24$ , the cost corresponding to the shortest path 1-3-5-6.
- Suppose that we choose to reduce the scalar  $\lambda_k$  by a factor of 2 whenever three successive iterations at a given value of  $\lambda_k$  have not improved on the best Lagrangian objective function value  $L(\mu)$ .
- The solution  $\mathbf{x}^0$  to the Lagrangian subproblem with  $\mu = \mathbf{0}$  corresponds to the path  $P = 1-2-4-6$ , the Lagrangian subproblem has an objective function value of  $L(\mathbf{0}) = 3$ , and the subgradient  $\mathcal{A}\mathbf{x}^0 - \mathbf{b}$  at  $\mu = \mathbf{0}$  is  $(t_p - 14) = 18 - 14 = 4$ .

# Subgradient Optimization and Inequality Constraints



- bold lines denote the shortest path with *Lagrange multiplier*  $\mu = 0$

# Subgradient Optimization and Inequality Constraints

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- So at the first step, we choose

$$\theta_k = \frac{\lambda_k[\text{UB} - L(\mu^k)]}{\|Ax^k - b\|^2}.$$

$$\theta_0 = 0.8(24 - 3)/16 = 1.05,$$

$$\mu^{k+1} = [\mu^k + \theta_k(Ax^k - b)]^+$$

$$\mu^1 = [0 + 1.05(4)]^+ = 4.2.$$

# Subgradient Optimization and Inequality Constraints

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- For  $\mu^1 = 4.2$ , the path  $P = 1-3-2-5-6$  solves the agrangian subproblem;
- Therefore,

$$L(\mu) = \min\{c_P + \mu(t_P - T) : P \in \mathcal{P}\}.$$

$$L(4.2) = 15 + 4.2(10) - 4.2(14) = 15 - 16.8 = -1.8,$$

- And  $\mathcal{A}x^1 - \mathbf{b}$  equals  $(t_p - 14) = 10 - 14 = -4$ .
- Since the path 1-3-2-5-6 is feasible, and its cost of 15 is less than **UB**, we change **UB** to value **15**.
- Therefore,

$$\theta_1 = 0.8(15 + 1.8)/16 = 0.84,$$

$$\mu^2 = [4.2 + 0.84(-4)]^+ = 0.84.$$

# Subgradient Optimization and Inequality Constraints

$k$	$\mu^k$	$t_p - T$	$L(\mu^k)$	$\lambda_k$	$\theta_k$
0	0.0000	4	3.0000	0.80000	1.0500
1	4.2000	-4	-1.8000	0.80000	0.8400
2	0.8400	4	6.3600	0.80000	0.4320
3	2.5680	-4	4.7280	0.80000	0.5136
4	0.5136	4	5.0544	0.80000	0.4973
5	2.5027	-4	4.9891	0.40000	0.2503
6	1.5016	1	6.5016	0.40000	3.3993
7	4.9010	-6	-5.4059	0.40000	0.2267
8	3.5406	-4	0.8376	0.40000	0.3541
9	2.1244	-4	6.5026	0.40000	0.2124
10	1.2746	1	6.2746	0.40000	3.4902
11	4.7648	-6	-4.5886	0.40000	0.2177
12	3.4589	-4	1.1646	0.20000	0.1729
13	2.7671	-4	3.9316	0.20000	0.1384
14	2.2137	-4	6.1453	0.20000	0.1107
15	1.7709	1	6.7709	0.20000	1.6458

# Subgradient Optimization and Inequality Constraints

16	3.4167	-4	1.3330	0.20000	0.1708
17	2.7334	-4	4.0664	0.20000	0.1367
18	2.1867	-4	6.2531	0.10000	0.0547
19	1.9680	1	6.9680	0.10000	0.8032
20	2.7712	-4	3.9150	0.10000	0.0693
21	2.4941	-4	5.0235	0.10000	0.0624
22	2.2447	-4	6.0212	0.05000	0.0281
23	2.1325	-4	6.4701	0.05000	0.0267
24	2.0258	-4	6.8966	0.05000	0.0253
25	1.9246	1	6.9246	0.00250	0.0202
26	1.9447	1	6.9447	0.00250	0.0201
27	1.9649	1	6.9649	0.00250	0.0201
28	1.9850	1	6.9850	0.00250	0.0200
29	2.0050	-4	6.9800	0.00250	0.0013
30	2.0000	-4	7.0000	0.00250	0.0012
31	1.9950	1	6.9950	0.00250	0.0200
32	2.0150	-4	6.9400	0.00250	0.0013
33	2.0100	-4	6.9601	0.00125	0.0006

## Subgradient Optimization and Inequality Constraints

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- From iterations 2 through 5, the shortest paths alternate between the paths 1-2-4-6 and 1-3-2-5-6.
- At the end of the fifth iteration, the algorithm has not improved upon (increased) the best Lagrangian objective function value of **6.36** for three iterations, so we reduce  $\lambda_k$  by a factor of 2.
- In the next 7 iterations the shortest paths are the paths 1-2-5-6, 1-3-5-6, 1-3-2-5-6, 1-3-2-5-6, 1-2-5-6, 1-3-5-6, and 1-3-2-5-6.

## Subgradient Optimization and Inequality Constraints

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- Once again for three consecutive iterations, the algorithm has not improved the best Lagrangian objective function value, so we decrease  $\lambda_k$  by a factor of 2 to value 0.2.
- From this point on, the algorithm chooses either path 1-3-2-5-6 or path 1-2-5-6 as the shortest path at each step.
- As we see, the Lagrangian objective function value is converging to the optimal value  $L^* = 7$  and the Lagrange multiplier is converging to its optimal value of  $\mu^* = 2$ .





The End